

CS16B Midterm 1

DSL = digital subscriber line
cable

enhanced → local → aggregate
switch

on-demand = statistical multiplexing better than reservations
for bursty

reservations - circuit switching = preconnection
on demand - packet switching - per packet, w/ buffer

delay

bandwidth = bits sent/received per time
prop delay = time to move through link x
bandwidth-delay product = bits in flight →
transmission delay = $\frac{\text{packet size}}{\text{transmission rate}}$
= packet size / transmission rate

propagation delay = link length / prop speed
queueing delay = time in buffer before processed
= $A \times W$ where A = length of queue packets in queue each sec
 W = waiting time
 A = arrival rate

processing delay = time to process packet

IP "best effort"
reliability @ TCP

forwarding = "dataplane" - directly use data packet & local forwarding info
routing = "control plane" - computing forwarding tables that guide packets, computed by routers
global forwarding table would if produces forwarding decision for data packets to destination
no dead ends, no loops

convergence delay - some link-state protocols - detect failure, flood with, recompute forwarding tables
last packets, keep by packets, out of order packets

loop-free routes = spanning tree avoid loops
Dijkstra's algorithm (Bellman-Ford - loop-free, exchange reach)
link state (OSPF) + distance vector (RIP) - loop-free routes = distance vector, pathing

intra-domain - link state (OSPF)
inter-domain - path vector (BGP)
convergence in minutes - take other way when connectivity broken
forwarded packet - if 2 routes through, 2nd route is first 2nd dest = ∞

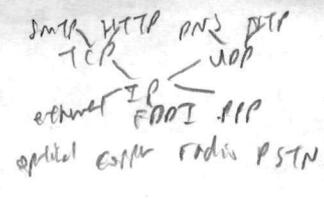
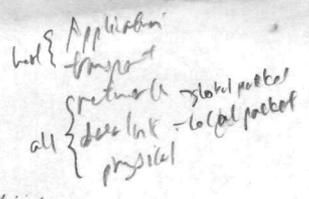
state = small forwarding tables at routers
churn = limited rate of change in configurations

IP address network 27 bits > 9
"multi-homed" problems (multiple providers)

egress - Customer → edge
peer → backbone
provider → backbone

AS PATH
LOCAL - PREFIX
MED

header: payload then fragment offset = address size
MTU
↑
header



IPsec = just a sec
IPsec - behind a tunnel

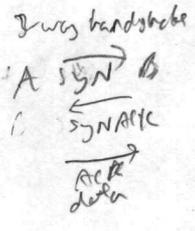


transport - communication between processes
application - addresses for application

	UDP	TCP
data	packets (header, payload)	stream of bytes of x length
state	best effort (IP)	reliable, in-order delivery, congestion control, flow control
	no ACKs, no seqs	ACKs, seqs, window

- Reliability
- checksums (error detection)
 - timers (loss detection)
 - ACKs (feedback)
 - exponential - all up to x
 - selective - X
 - sequence #s - duplicates, accuracy
 - sliding window efficiency

RTP: variable length, change in routing path
- first transmit
- next ACKs if order, hold receiver.
Reader



TCP buffer packet until all packets behind it have arrived + received or consumed or all bytes
"admitted window" = right hand edge of window must not exceed
- address of sending window.